

From Diophantus to Bitcoin...

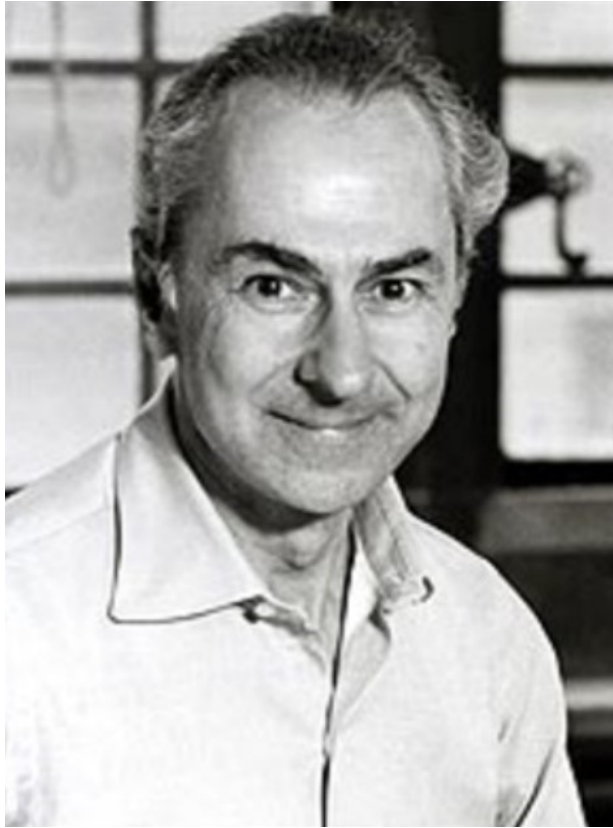


Why are elliptic curves everywhere?



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“It is possible to write endlessly
on elliptic curves.
(This is not a threat.)”

- Serge Lang
Diophantine Analysis.

PlayStation 3 hack - how it happened and what it means

(2010)

A group of coders claims the PS3 has been hacked, opening the doors to software piracy. We look into the implications



📷 The PS3 has been hacked. But how - and why? Photograph: Kevork Djansezian/AP

In December, a group of coders operating under the name [Fail0verflow](#) stood up at the Chaos Communications hackers conference in Berlin and proclaimed that the Sony PS3 security system was an epic fail. Through the use of what they termed "simple algebra" they had managed to exploit a weakness in the PlayStation 3's encryption system, thereby gaining the public key required to run any software on the machine.

BIZ & IT —

Google confirms critical Android crypto flaw used in \$5,700 Bitcoin heist

Java Crypto weakness could affect security in hundreds of thousands of apps.

DAN GOODIN - 8/14/2013, 9:15 PM

(2013)



$$e^{\pi\sqrt{163}} = 262537412640768743.\underbrace{999999999999}_{\text{twelve 9's!}}2500725\dots$$

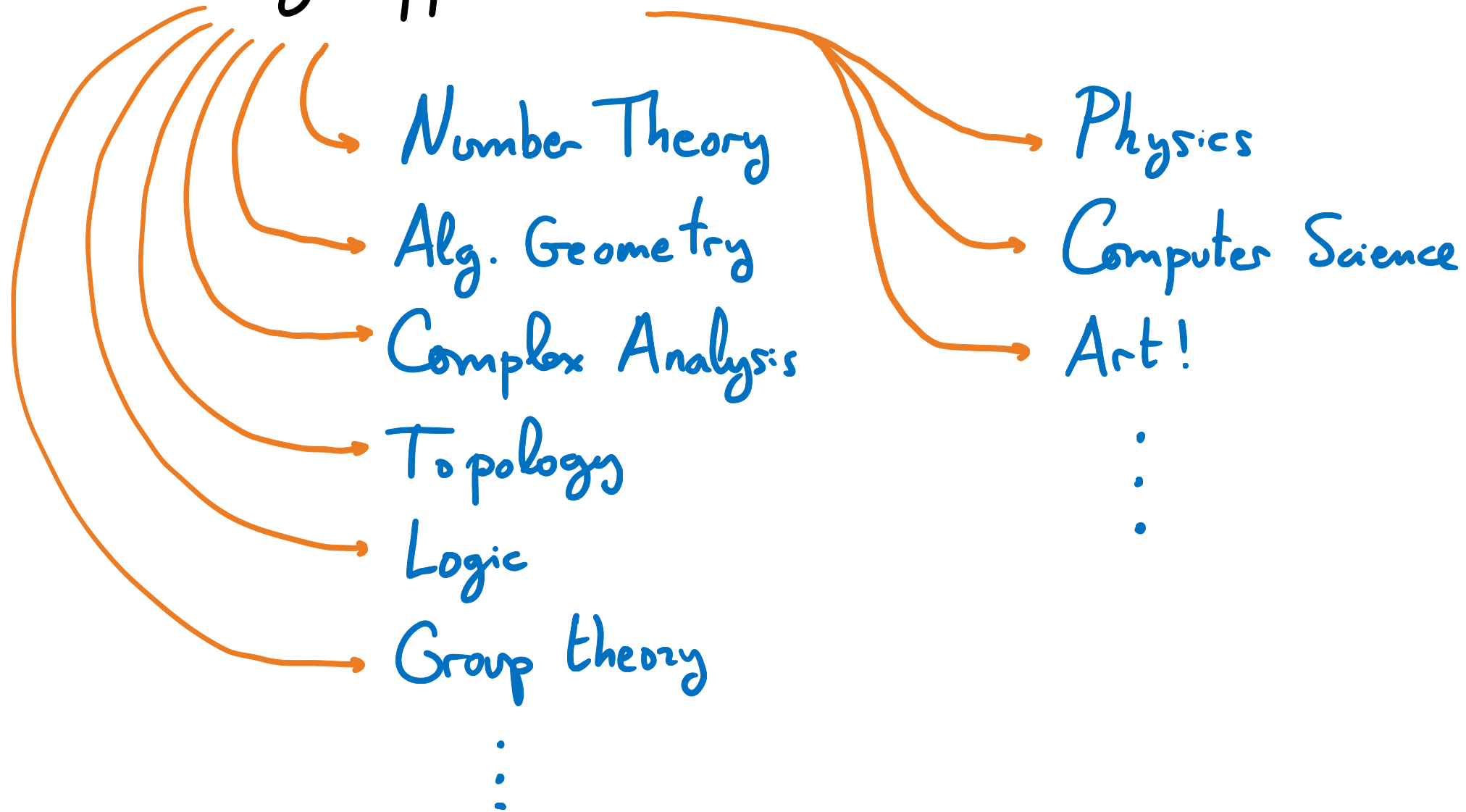
TRANSSCENDENTAL BY GELFOND - SCHNEIDER'S THEOREM!

$$p(n) = n^2 + n + 41$$

takes prime values for $n=0, 1, \dots, 39$

$$\begin{array}{ccccccccccc} 41, & 43, & 47, & 53, & \dots, & 1601, & 1681 = 41^2, & 1763 = 41 \cdot 43 \\ \uparrow & \uparrow & \uparrow & & & \uparrow & \uparrow & \uparrow \\ n=0 & n=1 & n=2 & \dots & n=39 & & n=40 & n=41 \end{array}$$

They appear in:



But ... Why?

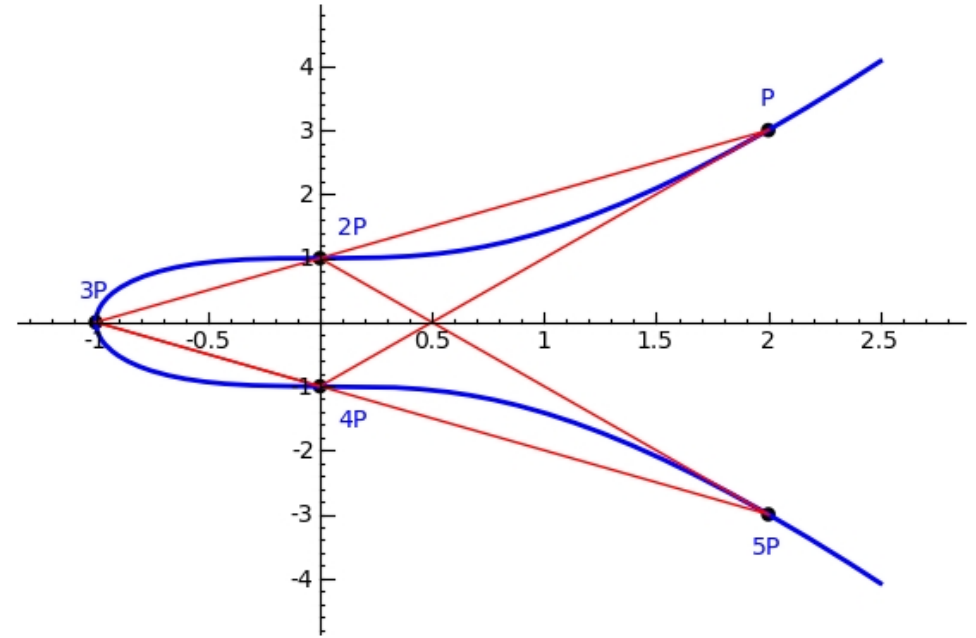
Because they are in the

"Sweet Spot"

of arithmetic, geometry, complex analysis,
complexity, and "cognition".

The Arithmetic "Sweet Spot"

What is an
elliptic curve?



What is a Diophantine Equation?

- $x^5 - x = 0$
- $3x + 5y = 1$
- $x^2 - 61y^2 = 1$
- $y^2 = x^3 - 157^2x$
- $x^3 + y^3 + z^3 = 42$



What is a Diophantine Equation?

More generally:

$$f(x_1, \dots, x_n) = 0$$

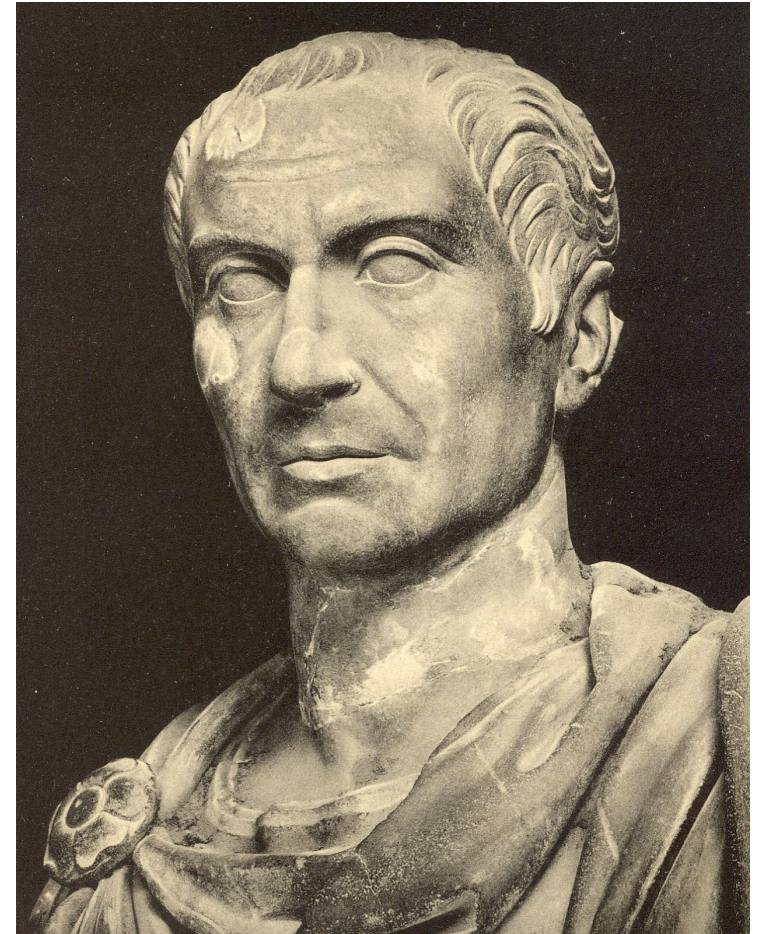
where $f \in \mathbb{Z}[x_1, \dots, x_n]$.



DIOPHANTUS OF ALEXANDRIA (BORN ~ 200 C.E.)

"The Father of Algebra"

- Known for the books
"Arithmetica,"
the first systematic study of
Diophantine equations.



DIOPHANTUS OF ALEXANDRIA (BORN ~ 200 C.E.)

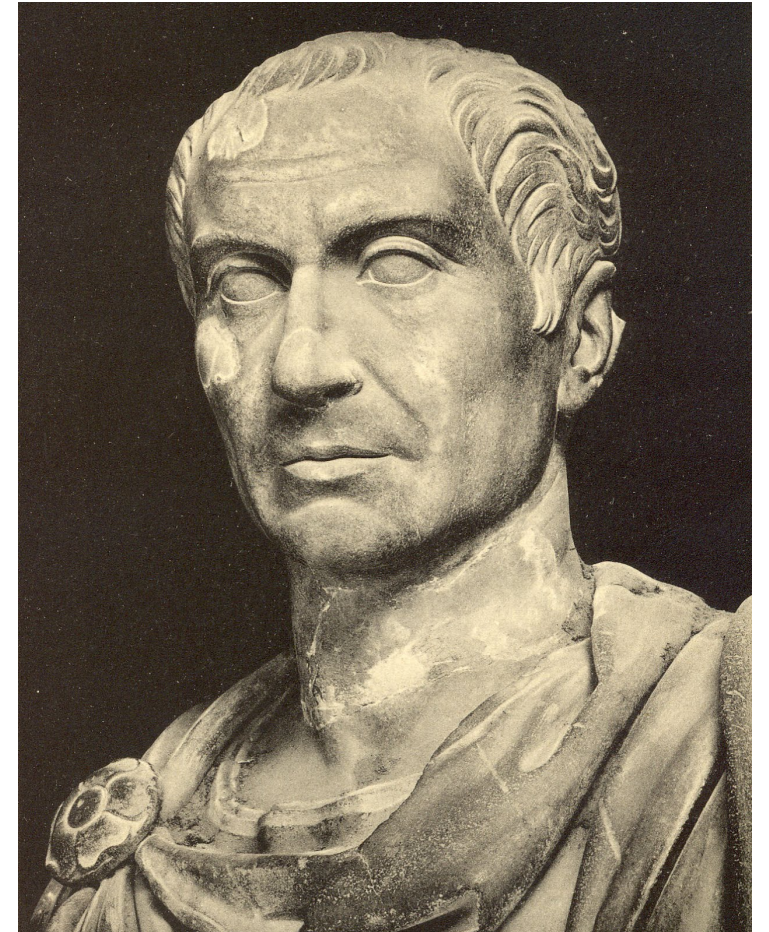
"The Father of Algebra"

DIOPHANTUS' EPITAPH:

'Here lies Diophantus,' the wonder behold.
Through art algebraic, the stone tells how old:

'God gave him his boyhood one-sixth of his life,
One twelfth more as youth while whiskers grew rife;
And then yet one-seventh ere marriage begun;
In five years there came a bouncing new son.
Alas, the dear child of master and sage
After attaining half the measure of his father's life chill fate took him.
After consoling his fate by the science of numbers for four years,
he ended his life.'

(METRODORUS, ~ 500 CE)

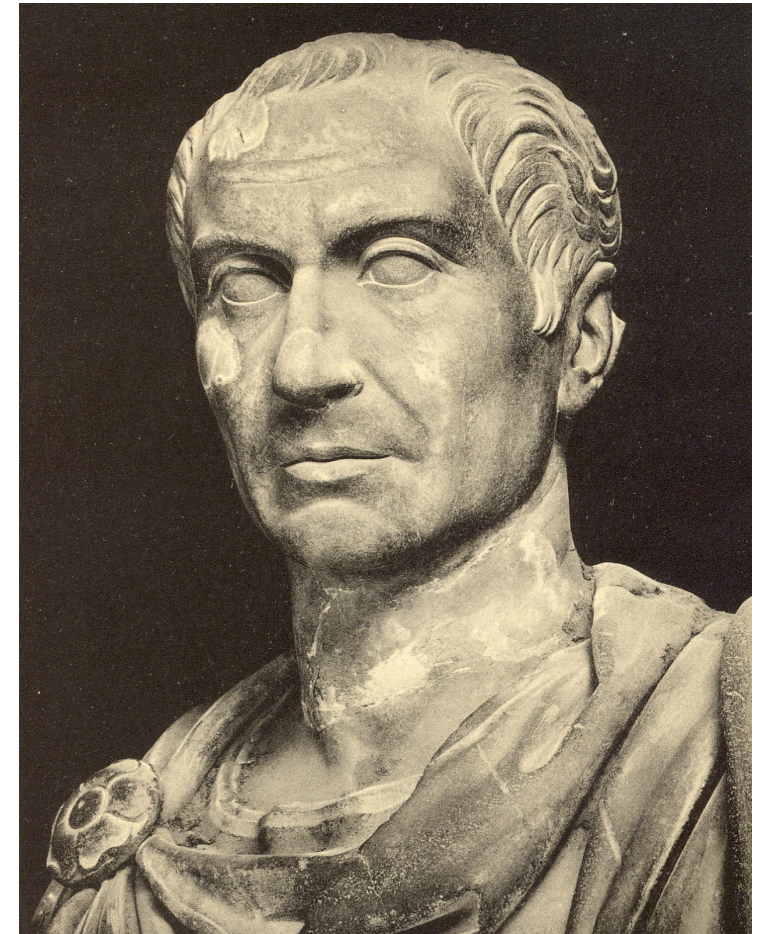


DIOPHANTUS OF ALEXANDRIA (BORN ~ 200 C.E.)

"The Father of Algebra"

BOOK IV, PROBLEM 24:

To divide a given number into two numbers such that their product is a cube minus its side.



DIOPHANTUS OF ALEXANDRIA (BORN ~ 200 C.E.)

"The Father of Algebra"

BOOK IV, PROBLEM 24:

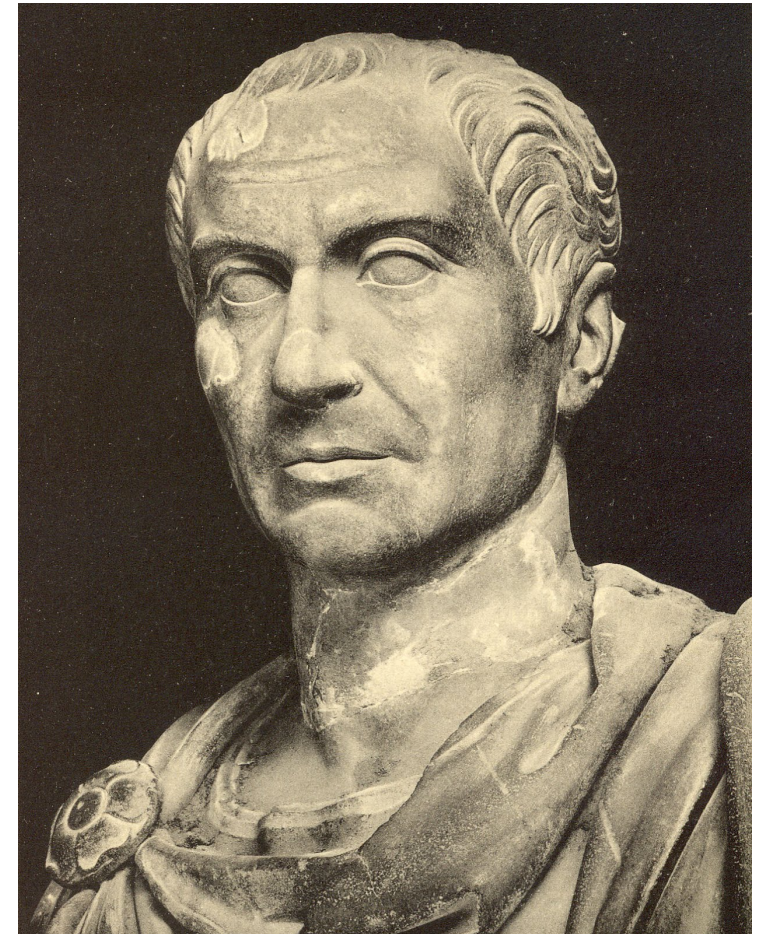
To divide a given number into two numbers such that their product is a cube minus its side.

$$a = y + (a-y)$$

such that

$$y \cdot (a-y) = x^3 - x$$

↖ an elliptic curve!



DIOPHANTUS OF ALEXANDRIA (BORN ~ 200 C.E.)

"The Father of Algebra"

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EXAMPLE

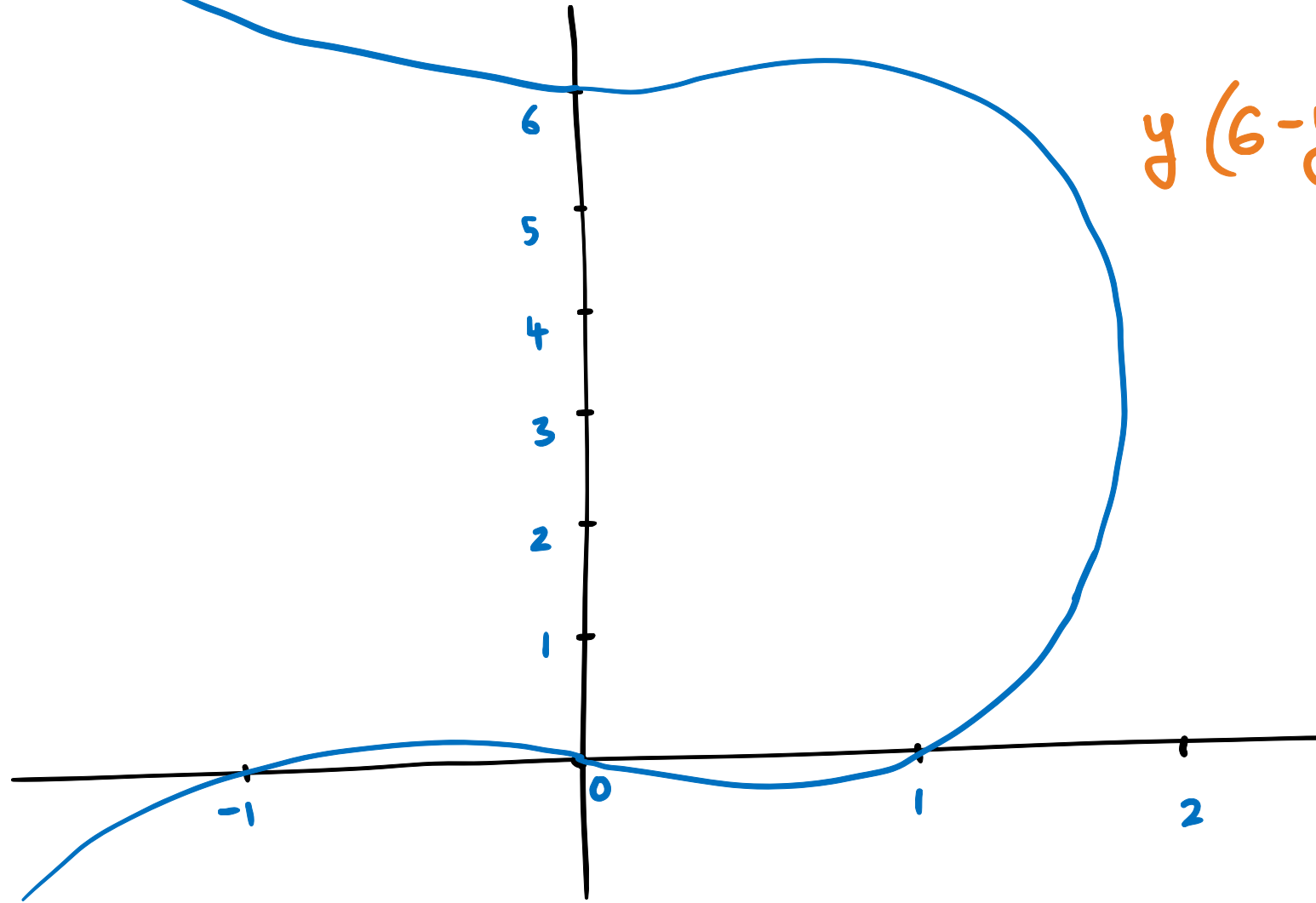
For $a = 6$

Diophantus finds:

$$y = \frac{26}{27}$$

How?

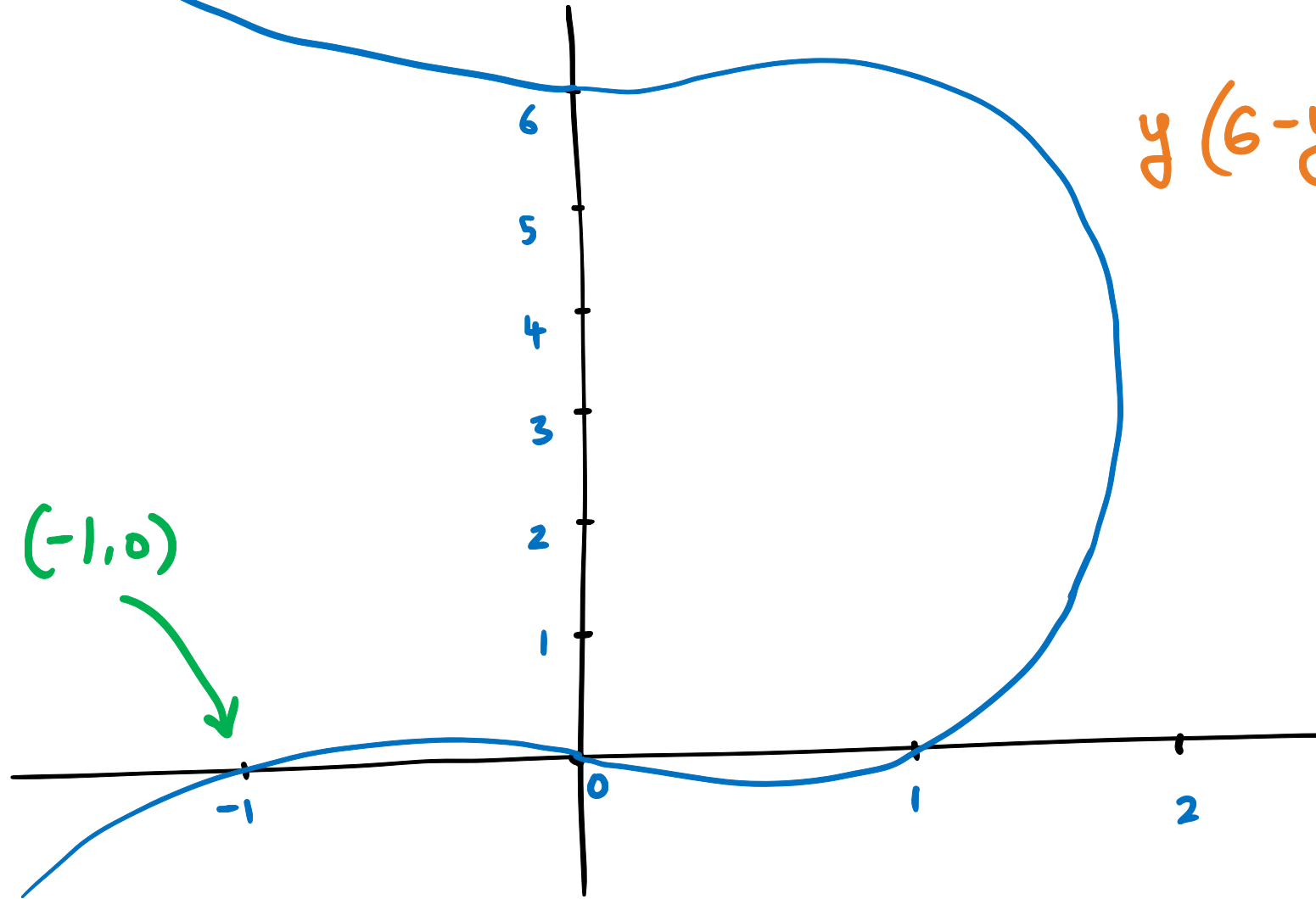
[Given a , find y s.t. $y(a-y) = x^3 - x$]



$$y(6-y) = x^3 - x$$

How?

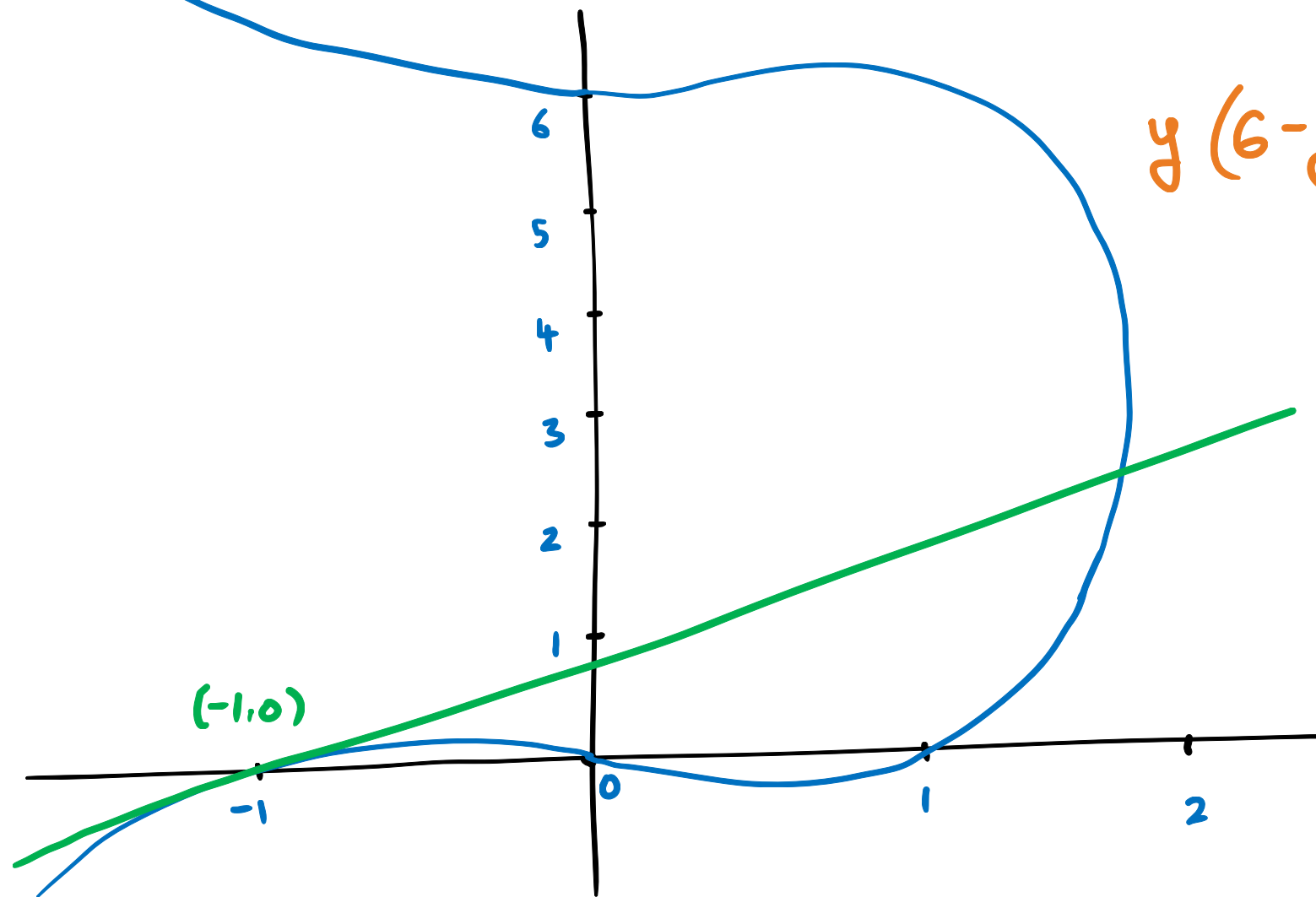
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$y(6-y) = x^3 - x$

How?

[Given a , find y s.t. $y(a-y) = x^3 - x$]



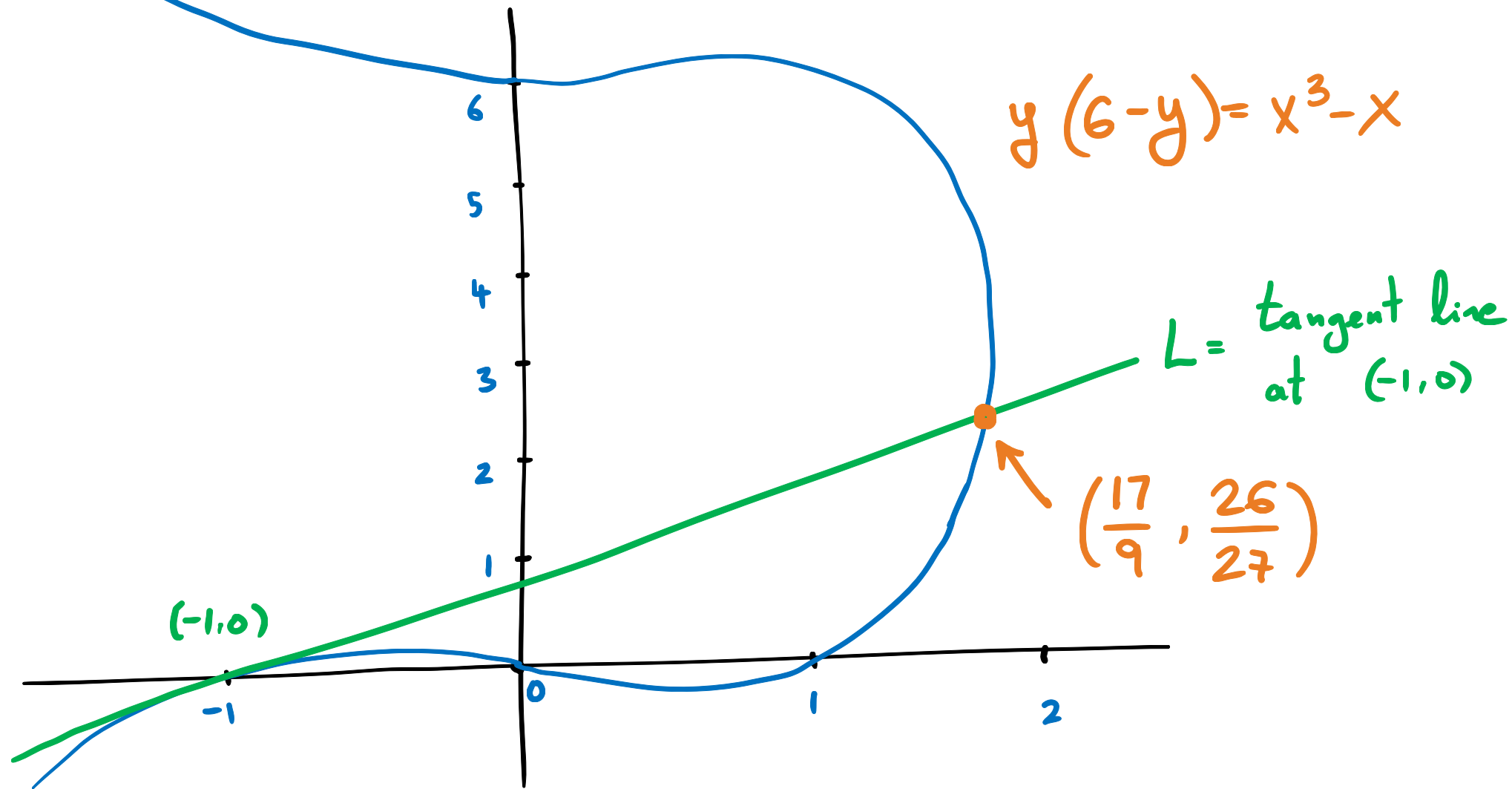
$y(6-y) = x^3 - x$

$L =$ tangent line at $(-1, 0)$

$(-1, 0)$

How?

[Given a , find y s.t. $y(a-y) = x^3 - x$]



$y(6-y) = x^3 - x$ is an example of an elliptic curve.

Diophantine Equations: $f(x_1, \dots, x_n) = 0$

deg f
number
of variables

• One variable: $f(x) = 0$ (a polynomial with \mathbb{Z} -coefficients)
($n=1$, arbitrary deg f)

• Two variables: $f(x, y) = 0$

(deg $f = 1$) Lines.

(deg $f = 2$) Conics and "products" of lines.

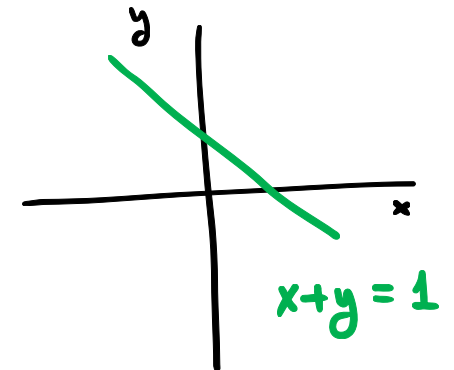
(deg $f = 3$) Smooth cubics and "products" of lines and conics.

⋮

• Two variables: $f(x,y)=0$. Can we find \mathbb{Z} or \mathbb{Q} solutions?

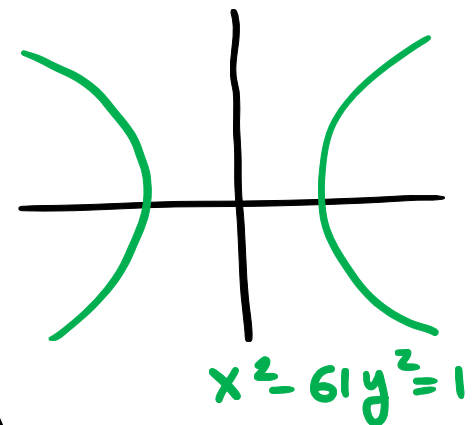
↳ (deg $f=1$) Lines.

↳ YES! Using divisibility and GCD's.



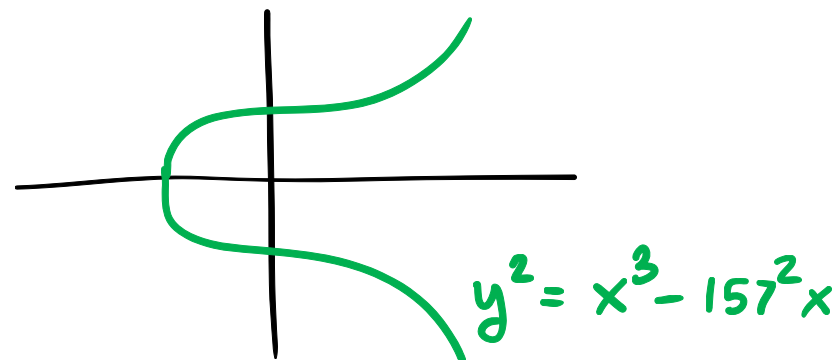
↳ (deg $f=2$) Conics and "products" of lines.

↳ YES! Hasse - Minkowski, Continued Fractions



(deg $f=3$) Smooth cubics and "products" of lines and conics.

↳ ???



ELLIPTIC CURVES

An elliptic curve over a field F is a smooth, projective curve of genus 1 with at least one point defined over F .

$$\begin{cases} y^2 = x^2(x-1) & \times \\ y^2 = x^3 & \sphericalangle \\ y = x^3 & \sphericalangle \end{cases}$$

NOT SMOOTH!

FACT: Every elliptic curve can be given by a Weierstrass model:

$$y^2 + a_1xy + a_3y = x^3 + a_2x^2 + a_4x + a_6, \quad a_i \in F.$$

If $\text{char } F \neq 2, 3$, then

$$y^2 = x^3 + Ax + B, \quad A, B \in F \\ 4A^3 + 27B^2 \neq 0.$$

INTEGER SOLUTIONS OF DIOPHANTINE EQUATIONS

... can be very intricate!

- $x^2 - 61 \cdot y^2 = 1$
 \searrow
 $\rightarrow 1^2 - 61 \cdot 0^2 = 1$
 $\rightarrow 1766319049^2 - 61 \cdot 226153980^2 = 1$

The solution:

$$1766319049^2 - 61 \cdot 226153980^2 = 1$$

comes from the continued fraction of $\sqrt{61}$:

$$\sqrt{61} = [7; \overline{1, 4, 3, 1, 2, 2, 1, 3, 4, 1, 14}]$$

$$= 7 + \frac{1}{1 + \frac{1}{4 + \frac{1}{3 + \frac{1}{1 + \frac{1}{2 + \frac{1}{2 + \frac{1}{1 + \dots}}}}}}}}$$

INTEGER SOLUTIONS OF DIOPHANTINE EQUATIONS

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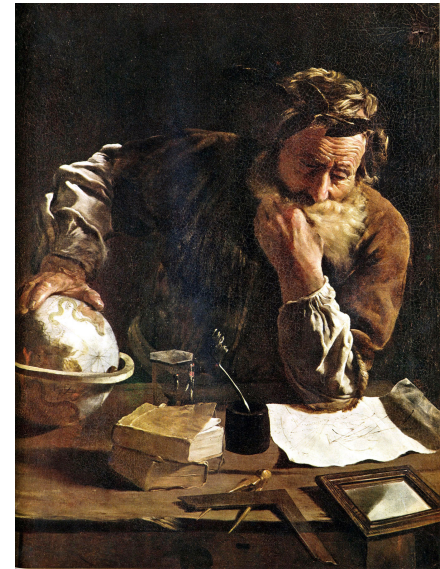
- Archimedes' cattle problem:

(287 - 212 BCE)

$$\left\{ \begin{array}{l} 5B = 6(W - Y) \\ 9D = 20(B - Y) \\ 13W = 42(D - Y) \\ 12W = 7(B + b) \\ 20b = 9(D + d) \\ 30d = 11(Y + y) \\ 42y = 13(W + w) \end{array} \right.$$

and

$$\left\{ \begin{array}{l} W + B = m^2 \\ 2(Y + D) = n(n + 1) \end{array} \right.$$



INTEGER SOLUTIONS OF DIOPHANTINE EQUATIONS

... can be very intricate!

- Archimedes' cattle problem:

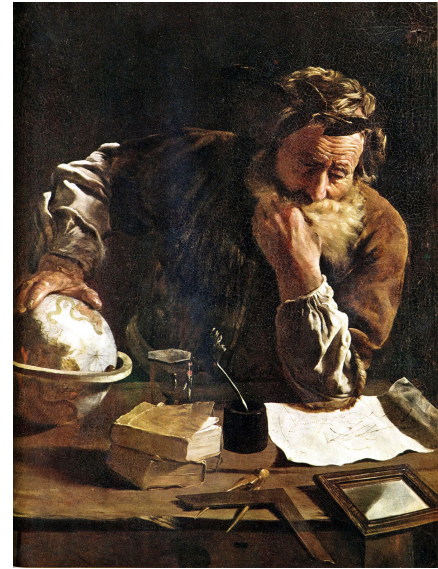
(287 - 212 BCE)

leads
to a Pell's
equation

$$u^2 - 4729494 \cdot v^2 = 1$$

leads
to a sol'n

$$7.76 \cdot 10^{206544} \text{ cattle.}$$



INTEGER SOLUTIONS OF DIOPHANTINE EQUATIONS

... can be very intricate!

- Is 42 a sum of three cubes?

$$X^3 + Y^3 + Z^3 = 42$$

The answer
to Life,
the Universe,
and Everything!
(Douglas Adams)

$$X = -80538738812075974$$

$$Y = 80435758145817515$$


$$Z = 12602123297335631$$

(Booker and
Sutherland, 2019)

Elliptic Curves in the Cognitive "Sweet Spot"

- Find all perfect squares that are one more than a cube.

Theory
of elliptic
curves


$$y^2 = x^3 + 1$$

(an elliptic curve!)


$$y^2 = 0, 1, \text{ or } 9.$$

Elliptic Curves in the Cognitive "Sweet Spot"

- Find all consecutive non-zero integers whose product is a perfect square:



$$y^2 = x(x+1)(x+2)$$

← elliptic curves

OR

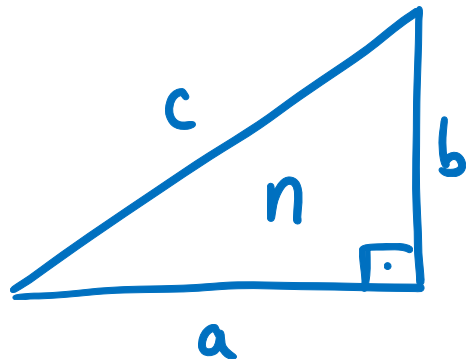
$$y^2 = (x-1)x(x+1) = x^3 - x$$

Theory of ell. curves

→ A: None.

Elliptic Curves in the Cognitive "Sweet Spot"

- Is there a right triangle with sides of rational length and...



- Area = 6 ? Yes! $(a, b, c) = (3, 4, 5)$
- Area = 5 ? Yes! $(a, b, c) = (\frac{20}{3}, \frac{3}{2}, \frac{41}{6})$ FIBONACCI 1200's
- Area = 1 ? No! Fermat 1600's
- Area = 157 ? Yes... Zagier 1900's

Elliptic Curves in the Cognitive "Sweet Spot"

- Is there a right triangle with sides of rational length and... area = n ?

↪ if and only if

$$y^2 = x^3 - n^2x \quad (\text{an elliptic curve!})$$

has a \mathbb{Q} -solution with $y \neq 0$.

Elliptic Curves in the Cognitive "Sweet Spot"

- Is there a right triangle with sides of rational length and... area = 157?

$$y^2 = x^3 - 157^2 x$$

leads to (a, b, c) with:

$$a = \frac{411340519227716149383203}{21666555693714761309610}, \quad b = \frac{6803298487826435051217540}{411340519227716149383203},$$

(due to D. Zagier)

$$c = \frac{224403517704336969924557513090674863160948472041}{8912332268928859588025535178967163570016480830}.$$

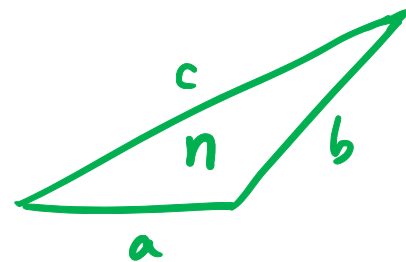
Elliptic Curves in the Cognitive "Sweet Spot"

- Is there a ~~right~~ triangle with sides of rational length and... area = n ?

Heron triangles (rational sides and area)

EXAMPLE: $n = 1$ is the area of

$$\left(\frac{3}{2}, \frac{5}{3}, \frac{17}{6}\right)$$



Elliptic Curves in the Cognitive "Sweet Spot"

- Is there a ~~right~~ triangle with sides of rational length and... area = n ?

Heron triangles (rational sides and area)

THEOREM (Goins, Maddox, 2006)

A number n is the area of a Heron triangle iff there is a rational number t s.t. $y^2 = x(x-nt)(x+nt)$ has a rational point with $y \neq 0$.

↖ an elliptic curve!

+ every n is the area of a Heron triangle.

Elliptic Curves in the Historical "Sweet Spot"

- Fermat's Last Theorem: $\begin{cases} x^n + y^n = z^n \\ xyz \neq 0 \end{cases}$ has no \mathbb{Z} -solutions for $n > 2$.

$x^3 + y^3 = z^3$ in $\mathbb{Z} \iff x^3 + y^3 = 1$ in \mathbb{Q} (an elliptic curve)

Hellegouarch, Frey: if $a^p + b^p = c^p$, study
 $E: y^2 = x(x - a^p)(x + b^p)$ (an elliptic curve??)

Serre, Ribet: E is semi-stable but not modular

Wiles, Taylor-Wiles: Every $\text{ell. curve}_{/\mathbb{Q}}$ is modular! ✓
↪ E cannot exist!

INTEGER SOLUTIONS OF DIOPHANTINE EQUATIONS

... can be very intricate! Algorithm?

Hilbert's Tenth Problem (1900)

Given a Diophantine equation with any number of unknown quantities and with rational integral numerical coefficients:

To devise a process according to which it can be determined in a finite number of operations whether the equation is solvable in rational integers.



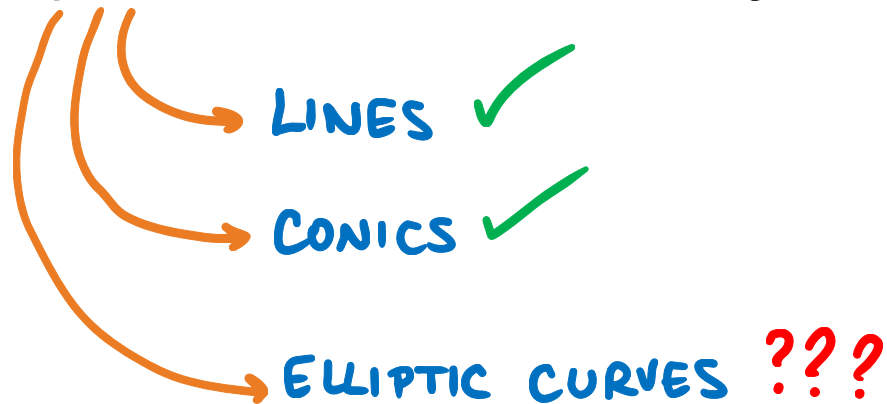
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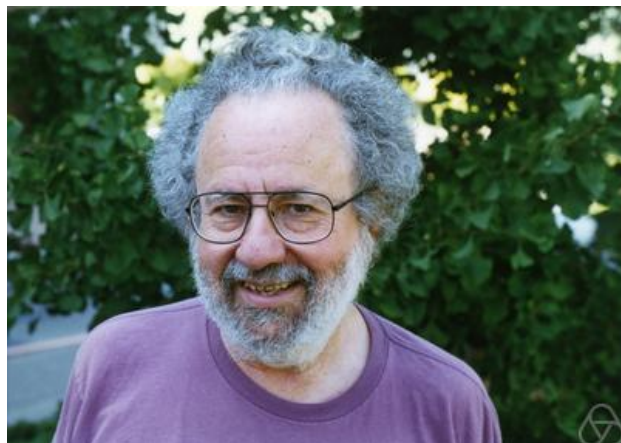
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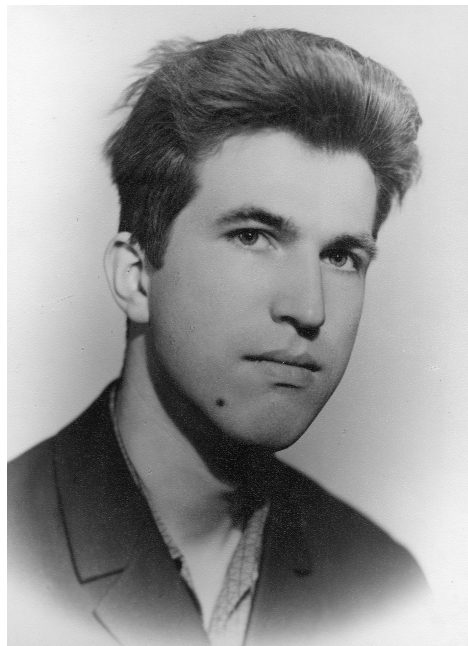
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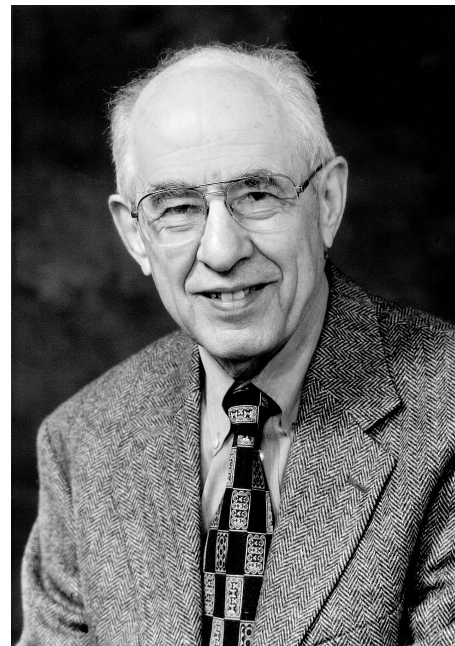
XXTH CENTURY SURPRISE ^ LOGIC!



DAVIS



MATIYASEVICH



PUTNAM



ROBINSON

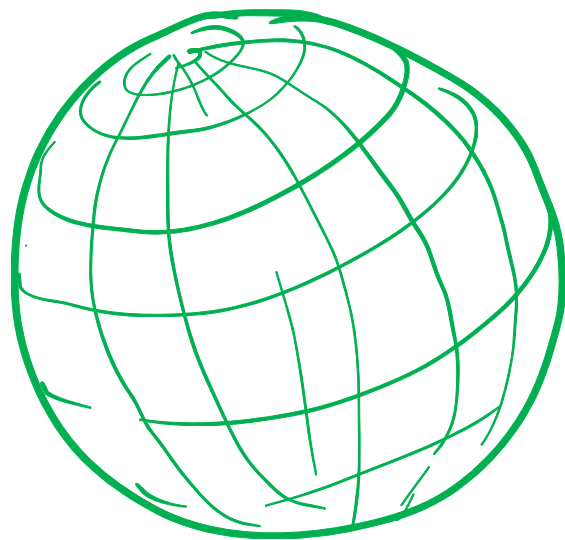
HILBERT'S 10TH (OVER \mathbb{Z}) : **NO!** THERE IS NO SUCH ALGORITHM.
(1970)

The Geometric "Sweet Spot"

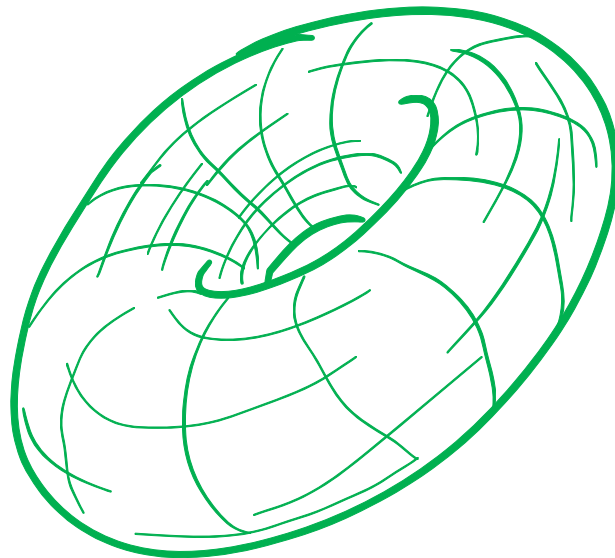
- Complex algebraic curves are classified by their genus:

$C : f(x, y) = 0 \quad f \in \mathbb{Z}[x, y].$
over \mathbb{C} classified according to their genus

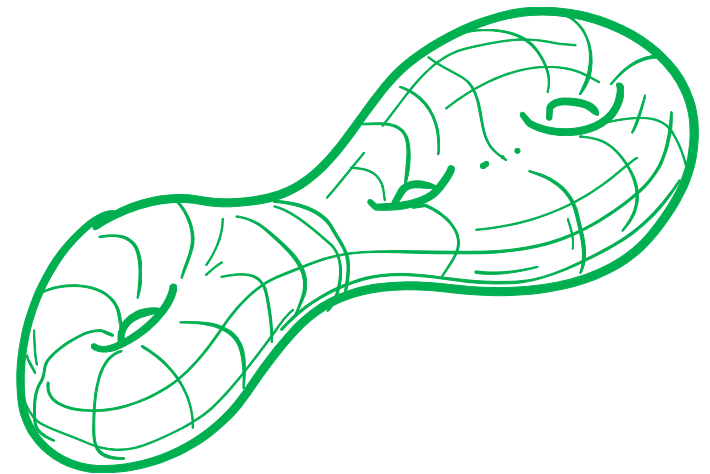
dim. of space of regular diff. 1-forms on C .



$g = 0$



$g = 1$



$g \geq 2$

The Geometric "Sweet Spot"

- Complex algebraic curves are classified by their genus:

$$C : f(x, y) = 0 \quad f \in \mathbb{Z}[x, y].$$

over \mathbb{C} classified according to their genus

dim. of space of regular diff. 1-forms on C .

<u>Curves</u>	<u>deg f</u>	<u>genus</u>	<u>#rational points over \mathbb{Q}</u>
Lines, Conics	1, 2	0	\emptyset or ∞ 'ly many
Cubics ⁺	3, 4	1	\emptyset , or finitely many, or ∞ 'ly many
Higher genus	≥ 4	≥ 2	\emptyset or finitely many. (Faltings')

EXAMPLES

The Geometric "Sweet Spot"

$$g=0 \begin{cases} \bullet x^2 + y^2 = -1 & \text{no } \mathbb{Q}\text{-points} \\ \bullet x^2 + y^2 = 1 & \infty\text{'ly many : } \left\{ \left(\frac{t^2-1}{t^2+1}, \frac{2t}{t^2+1} \right) : t \in \mathbb{Q} \right\} \cup \{(1,0)\} \end{cases}$$

$$g=1 \begin{cases} \bullet 3x^3 + 4y^3 + 5 = 0 & \text{no } \mathbb{Q}\text{-points (one pt. at } \infty : [0,1,0]) \\ \bullet x^3 + y^3 = 1 & 2 \mathbb{Q}\text{-points (plus one pt. at } \infty) \\ \bullet y^2 = x^3 - 2 & \infty\text{'ly many } \mathbb{Q}\text{-points} \end{cases}$$

$y^2 - 9y = x^3 - 27$

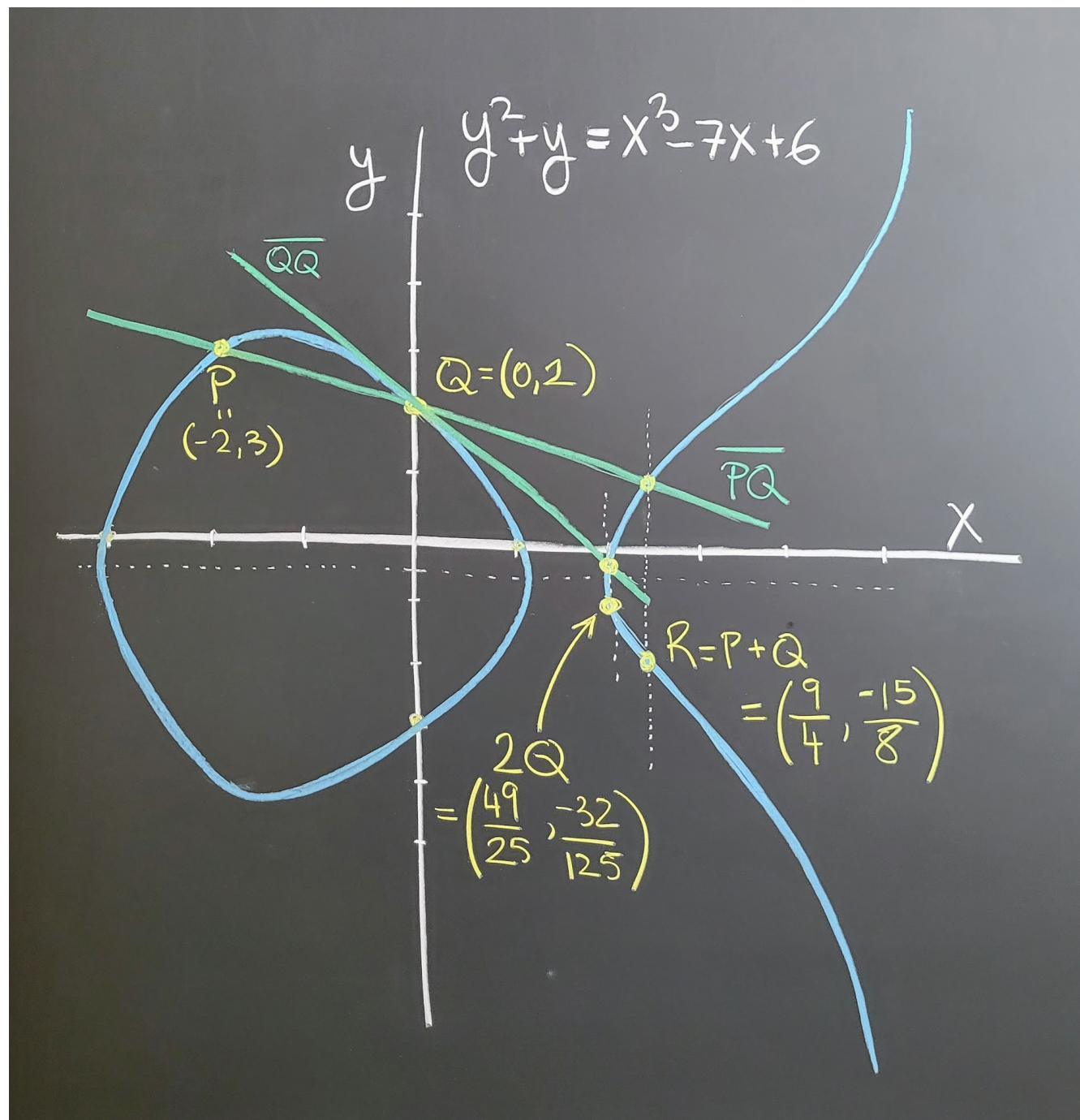
$$g \geq 2 \begin{cases} \bullet_3 x^4 + y^4 = -1 & \text{no } \mathbb{Q}\text{-points} \\ \bullet_2 y^2 = x - x^5 & 3 \mathbb{Q}\text{-points (plus one at } \infty) \end{cases}$$

Moreover!

$g = 1$
+
one pt.

} \Rightarrow

Geometric
Group
Structure!



Mordell (-Weil), 1922

Let E/\mathbb{Q} be an elliptic curve. Then, the set of rational points on E , denoted by $E(\mathbb{Q})$, is a finitely generated abelian group.

$$E(\mathbb{Q}) = E(\mathbb{Q})_{\text{tors}} \oplus \mathbb{Z}^{R_{E/\mathbb{Q}}}$$



Mazur:
15 possibilities



Elkies:
 $\sup \{R_{E/\mathbb{Q}}\} \geq 28.$

RANKS

Logic "Sweet Spot"

- Hilbert's 10th Problem ...
... over other rings?

EXAMPLE: Replace \mathbb{Z} by $\mathbb{Z}[i] = \{a+bi : a, b \in \mathbb{Z}, i^2 = -1\}$

- Is there an algorithm to determine (in a finite number of steps) if a diophantine equation over $\mathbb{Z}[i]$ has a solution over $\mathbb{Z}[i]$?

Logic "Sweet Spot"

- Hilbert's 10th Problem ...
... over other rings?

Denef, Pheidas, Shlapentokh, Poonen:

Theorem Let F be a number field. If there is an elliptic curve E/\mathbb{Q} s.t.

$$R_{E/F} = R_{E/\mathbb{Q}} = 1$$

then H10 is false over \mathcal{O}_F , the ring of integers of F .

Example

$$y^2 + y = x^3 - x$$

$$R_{E/\mathbb{Q}(i)} = R_{E/\mathbb{Q}} = 1$$

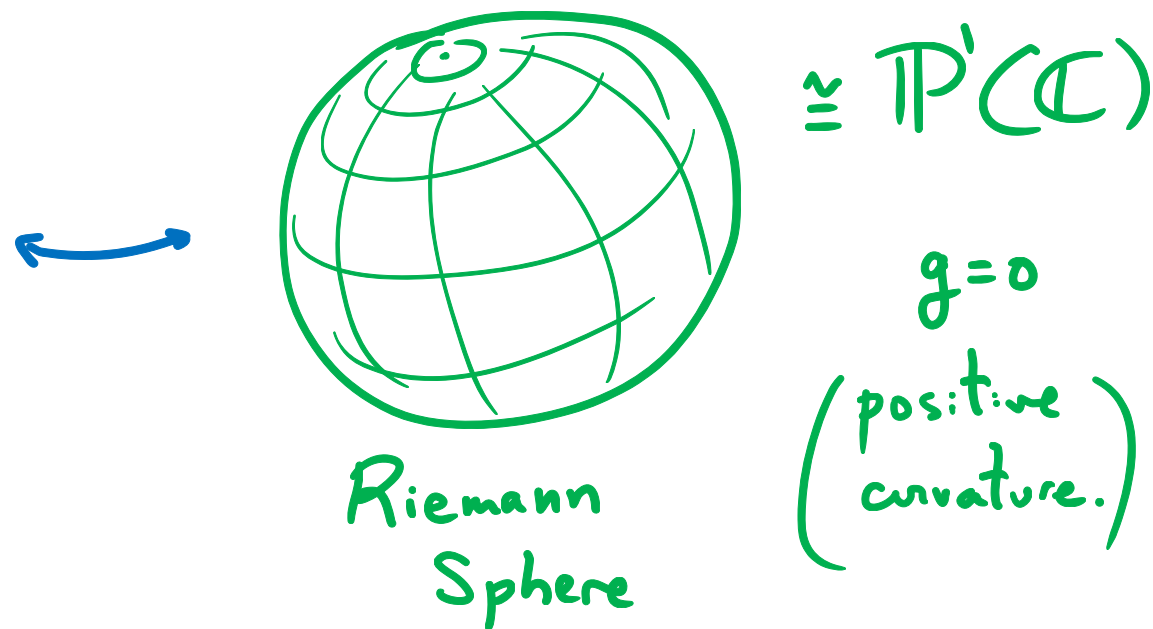
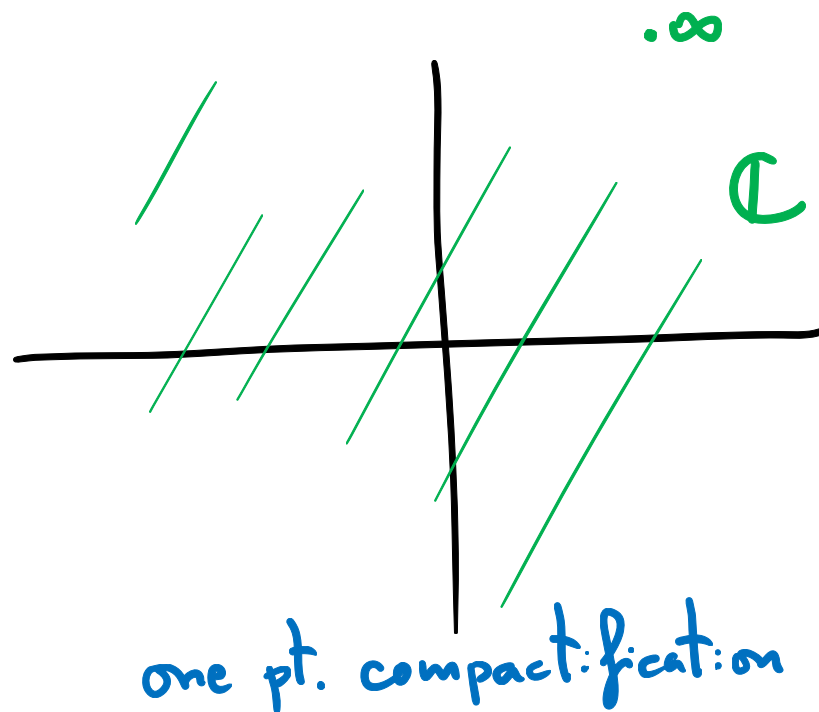
\Rightarrow H10/ $\mathbb{Z}[i]$
no algorithm exists!

Complex Analysis: "Sweet Spot"



A central object of study is a Riemann Surface:
a connected one-dimensional complex manifold.

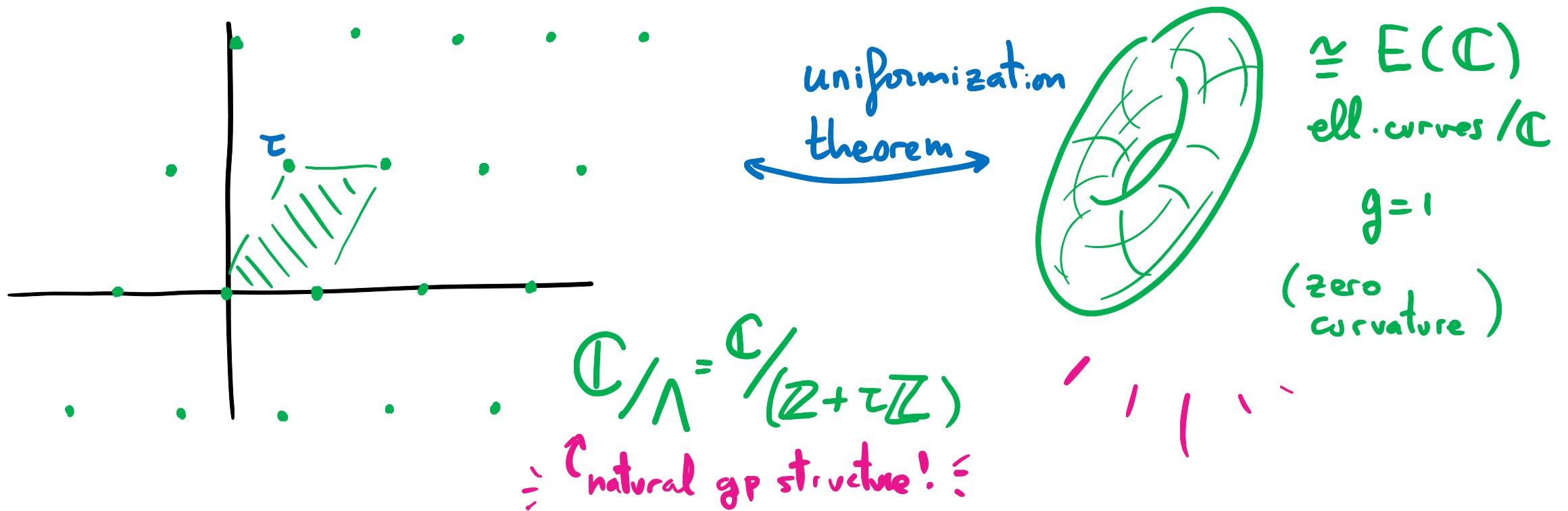
FACT: $\{\text{compact R.S.}\} \longleftrightarrow \{\text{complex smooth proj. alg. curves}/\mathbb{C}\}$



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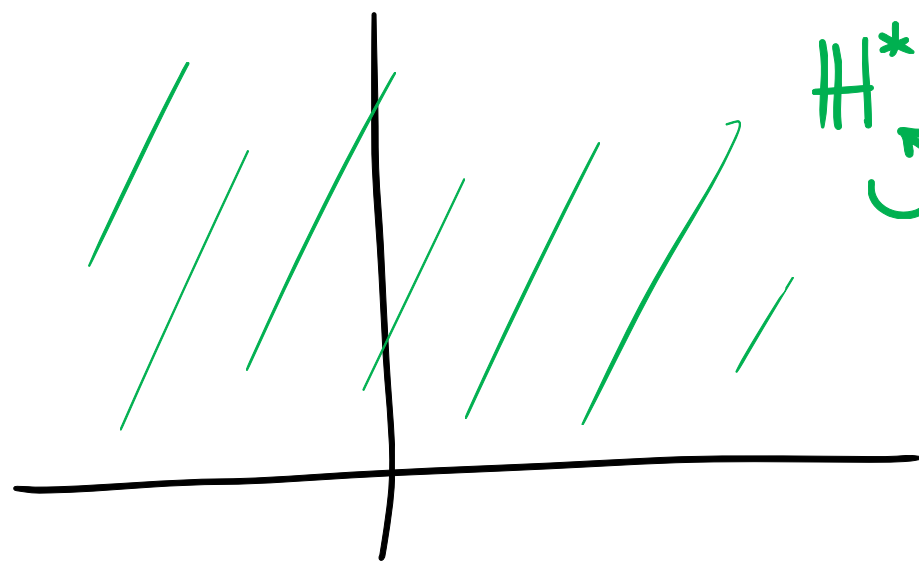
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Complex Analysis: "Sweet Spot"

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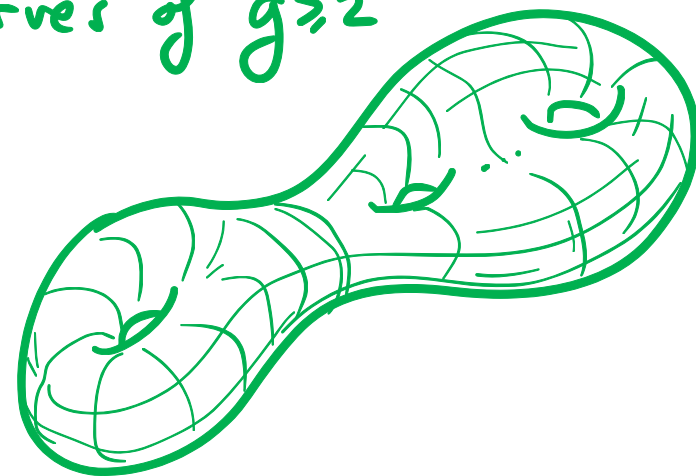


H^*
 $\curvearrowright \Gamma$

H^*
 Γ
(neg. curvature)



Curves of $g \geq 2$



\mathbb{H}^*

hyperbolic
Poincare
plane

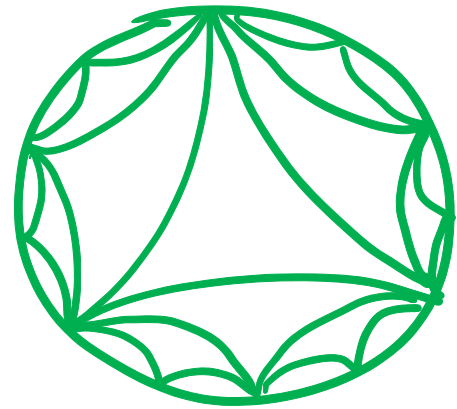
Holomorphic
Diff. Forms

Modular
Forms

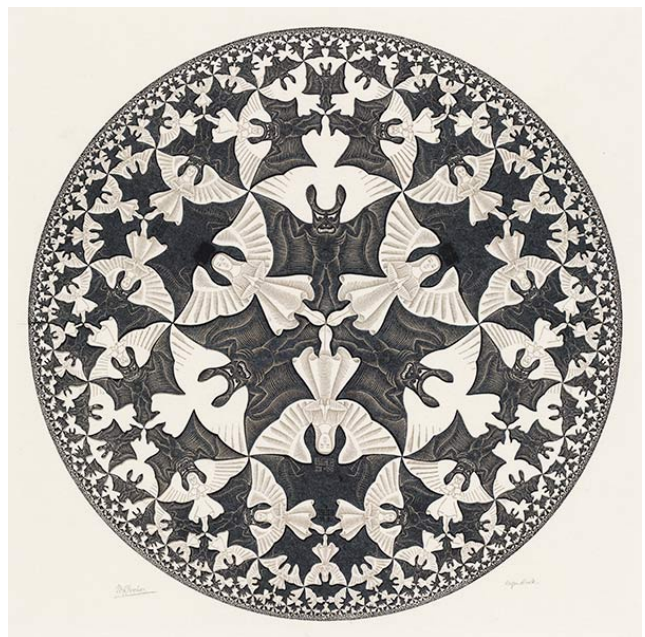
Elliptic
Curves!

Algebraic
Topology

Hopkins - Mahowald - Miller
Topological
Modular Forms



Art!



Escher
Tilings.



Complex Analysis: "Sweet Spot"

$$\{\mathbb{C}/\Lambda_\tau\}_{/iso} \longleftrightarrow \{E/\mathbb{C}\}_{/iso} \longleftrightarrow \mathbb{C}$$

$$\begin{array}{ccc} \mathbb{C}/\Lambda_\tau & \longleftrightarrow & y^2 = 4x^3 - g_2(\Lambda)x - g_3(\Lambda) \longleftrightarrow j(\tau) \\ z \bmod \Lambda_\tau & \longleftrightarrow & (f(z, \Lambda), f'(z, \Lambda)) \quad \frac{12^3 \cdot g_2(\Lambda)^3}{g_2(\Lambda)^3 - 27g_3(\Lambda)^2} \end{array}$$

where

$$j(\tau) = \frac{1}{q} + 744 + 196884q + 21493760q^2 + \dots \quad \text{w/ } q = e^{2\pi i \tau}$$

is the modular j -invariant function.

Group Theory ??

$$j(\tau) = \frac{1}{q} + 744 + \underbrace{196884 q}_{\Gamma_1 + \Gamma_2} + \underbrace{21493760 q^2}_{\Gamma_1 + \Gamma_2 + \Gamma_3} + \dots \quad \text{w/ } q = e^{2\pi i \tau}$$

$\Gamma_1 = 1$

$\Gamma_1 + \Gamma_2$

$\Gamma_1 + \Gamma_2 + \Gamma_3$

(John McKay)
in 1978

Group Theory ??

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$\Gamma_1 = 1$ $\Gamma_1 + \Gamma_2$ $\Gamma_1 + \Gamma_2 + \Gamma_3$

where $\Gamma_i = \text{dim. of irreducible representations of } M$, the **MONSTER GROUP!**

(John McKay)
in 1978

↓
Conway-Norton

↓
Borcherds, 1992 ✓

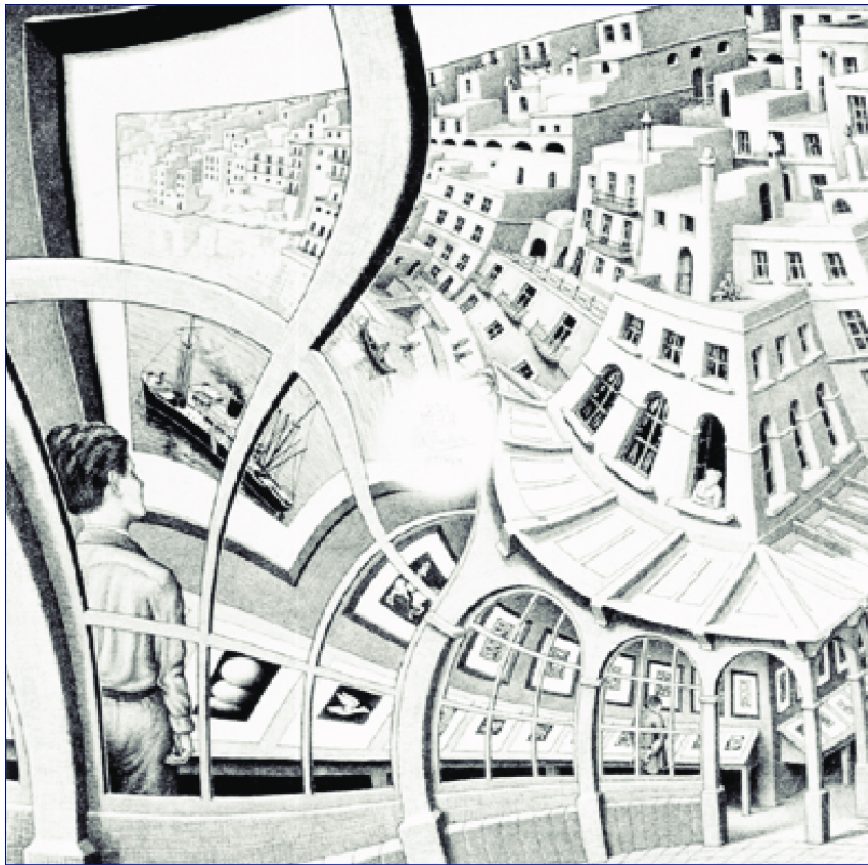
$$j(\tau) = \frac{1}{q} + 744 + 196884q + 21493760q^2 + \dots \quad \text{w/ } q = e^{2\pi i \tau}$$

$$E: y^2 + y = x^3 - 2174420x + 1234136692$$

$\tau = \frac{1 + \sqrt{-163}}{2} + CM \Rightarrow j(\tau) \text{ is an integer}$

$$-\frac{1}{q} = e^{\pi \sqrt{163}} \text{ is very close to an integer}$$

C/A and Art!



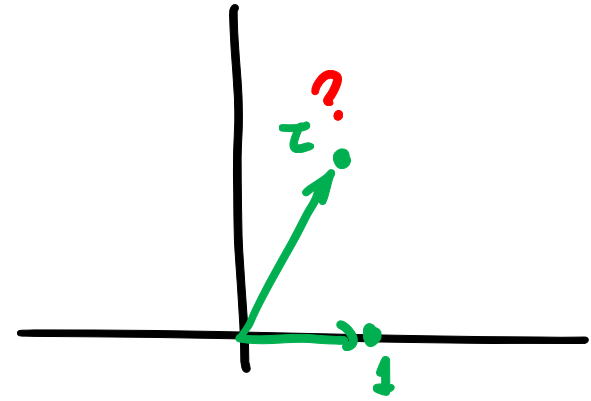
Escher's incomplete
"Prentententoonstelling"

→
de Smit
+
Lenstra

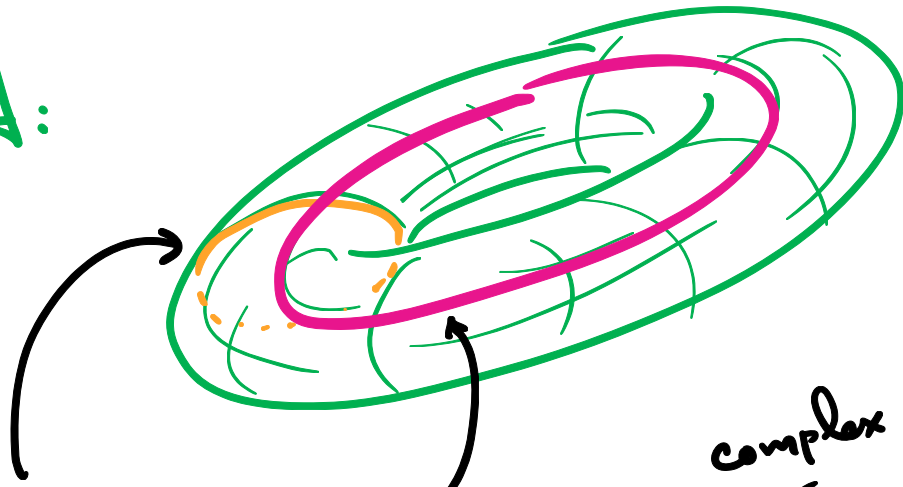


Escher's complete
"Prentententoonstelling"

Q Given E/\mathbb{C} what is Λ_E ?



A:



two line integrals give γ ^{complex values} w_1, w_2 s.t. $\Lambda = \langle w_1, w_2 \rangle$

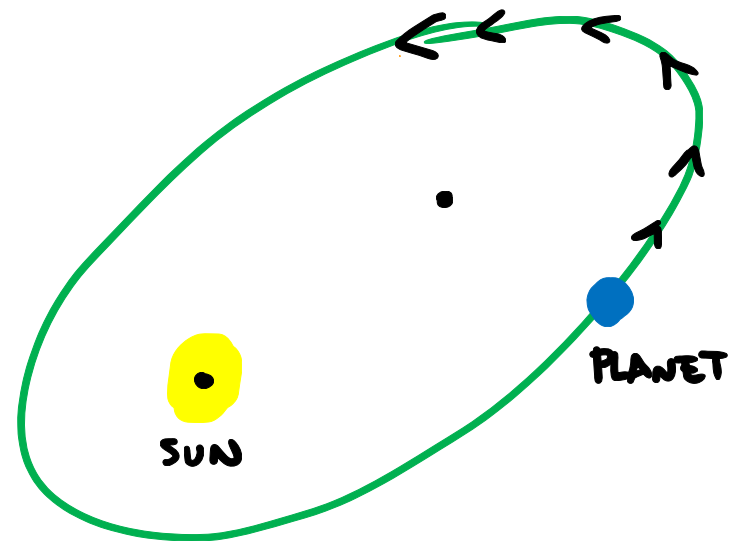
Why "elliptic" curves?

Elliptic curves in the Physics "Sweet Spot"

... in the 1600's!

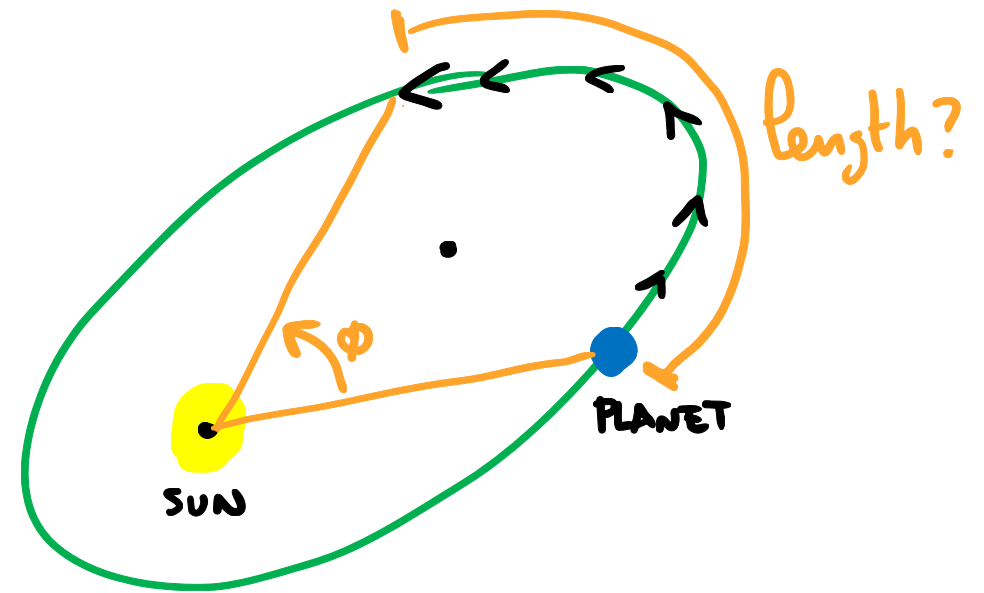
Kepler's Laws (1609-1619):

Planets move in elliptical orbits around the Sun.



Kepler's Laws (1609-1619):

Planets move in elliptical orbits around the Sun.



Q Arc length on an ellipse?

Sol'n as an inf. series: Newton, Euler,...

As an integral?

$$\frac{x^2}{a^2} + \frac{y^2}{b^2} = 1$$

Circumf.: $4 \int_0^{\pi/2} \sqrt{a^2 \cos^2 t + b^2 \sin^2 t} dt = 2 \int_{b^2/a^2}^1 \frac{t dt}{\sqrt{t(t-1)(t-b^2/a^2)}}$

Arc length: $u = F(\phi) \longleftrightarrow \phi = F^{-1}(u)$

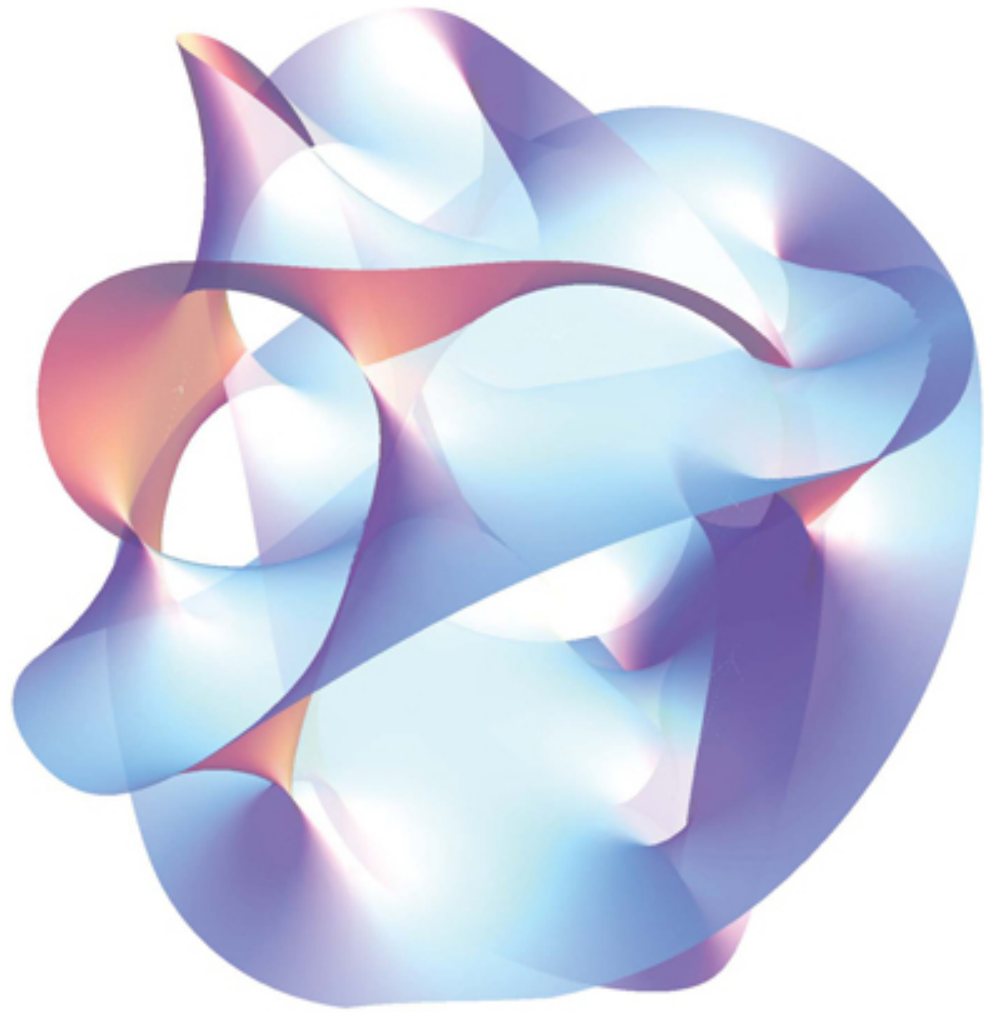
Jacobi
Abel

ELLIPTIC
FUNCTIONS

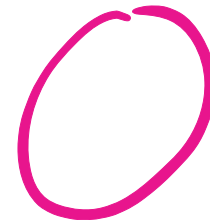
$f(z)$

$$y^2 = x(x-1)(x-b^2/a^2)$$

Physics in the 1900's : String Theory



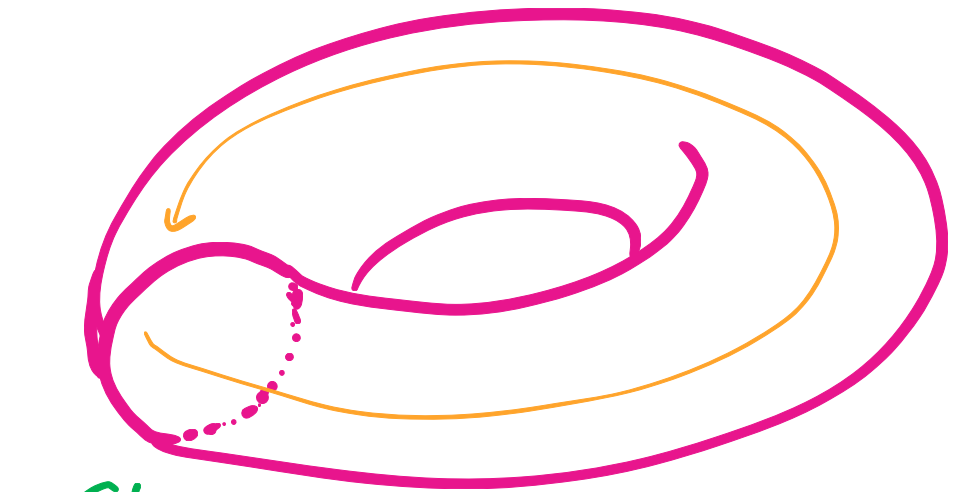
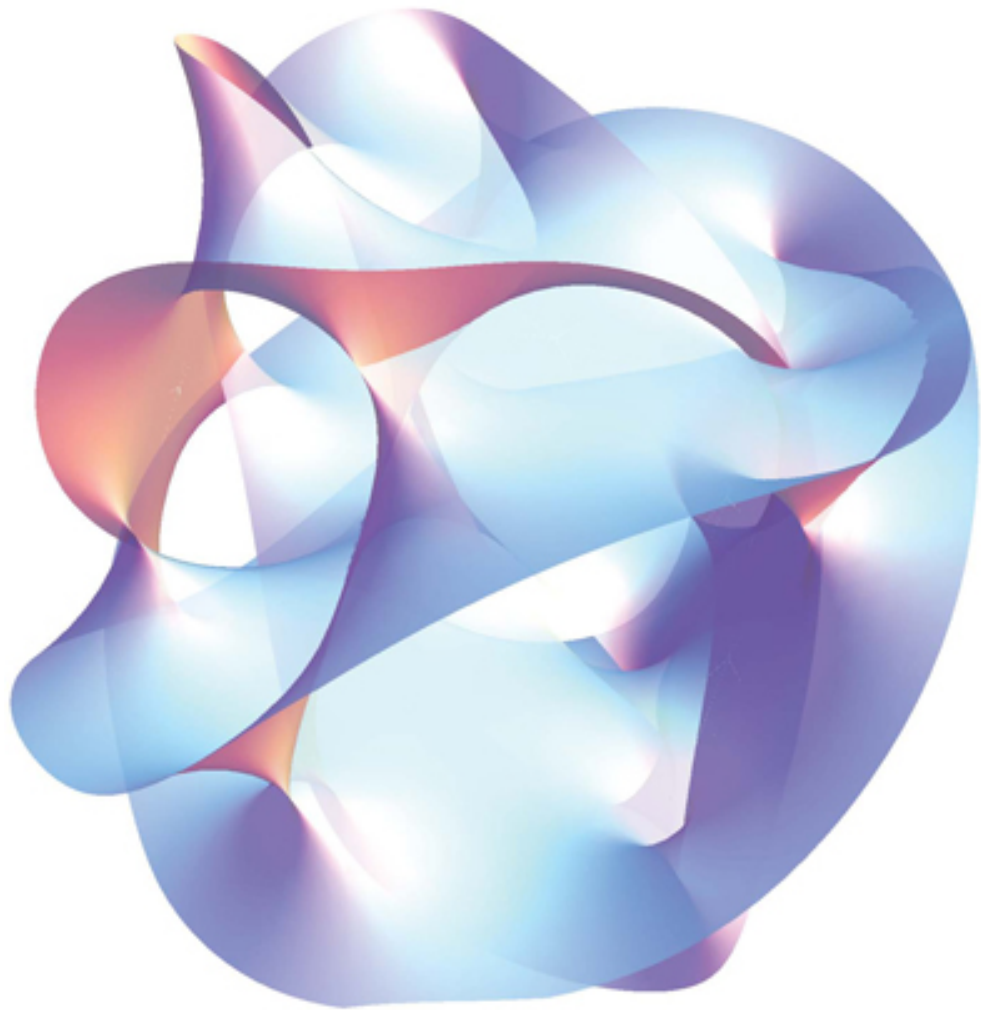
Space time in 6- or 23-dimensional
Calabi-Yau manifold



String

Physics in the 1900's : String Theory

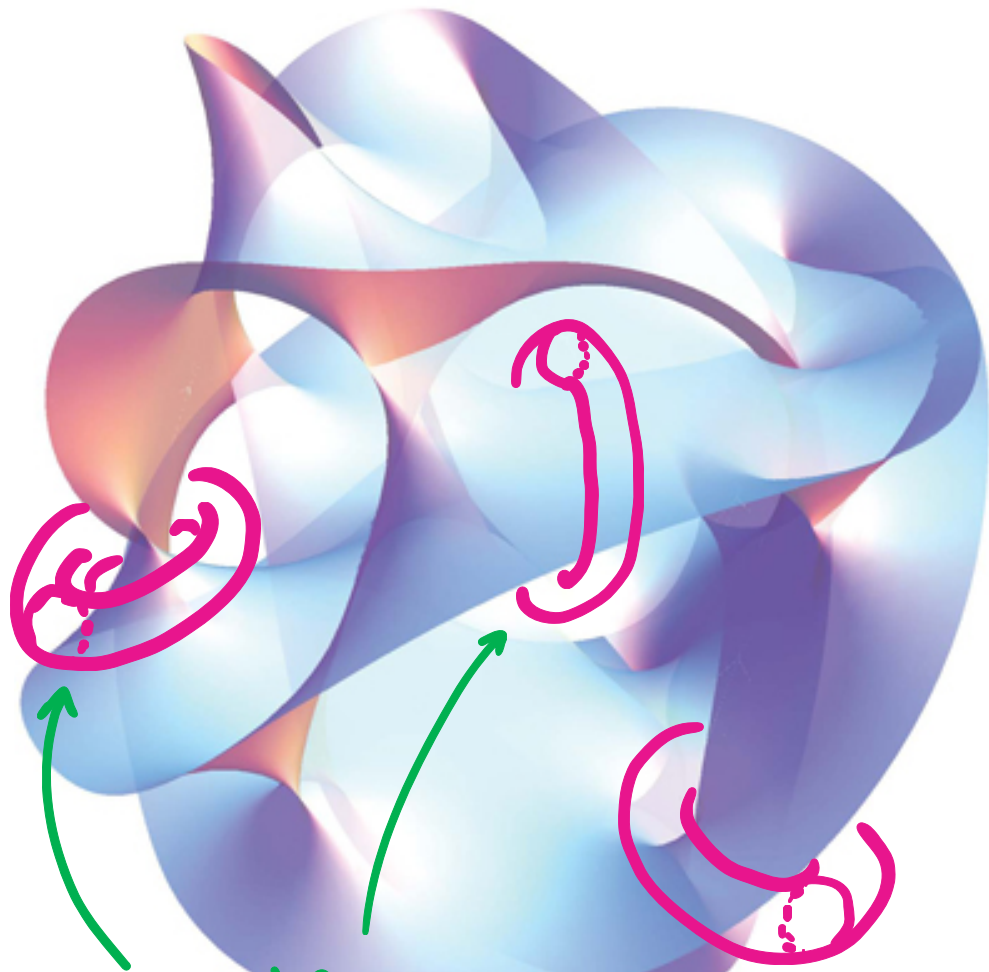
Space time in 6- or 23-dimensional
Calabi - Yau manifold



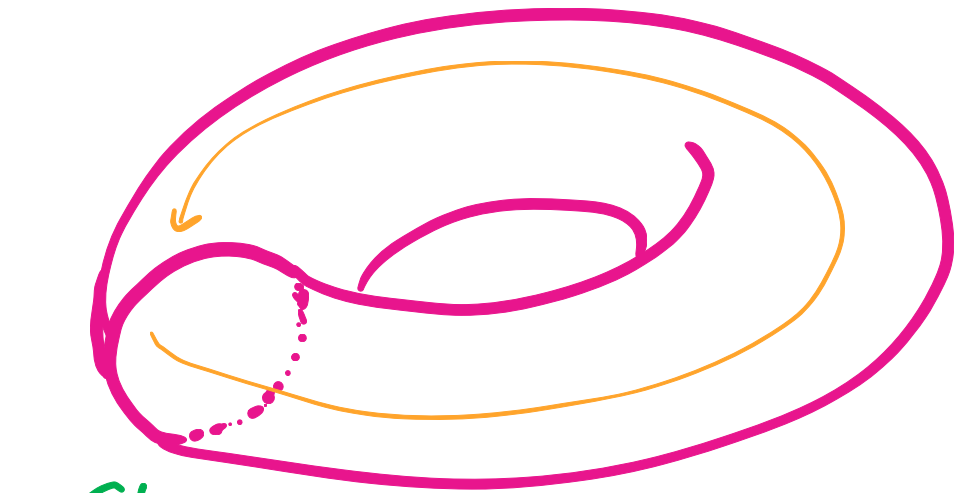
String
Closed Path

Physics in the 1900's : String Theory

Space time in 6- or 23-dimensional
Calabi-Yau manifold



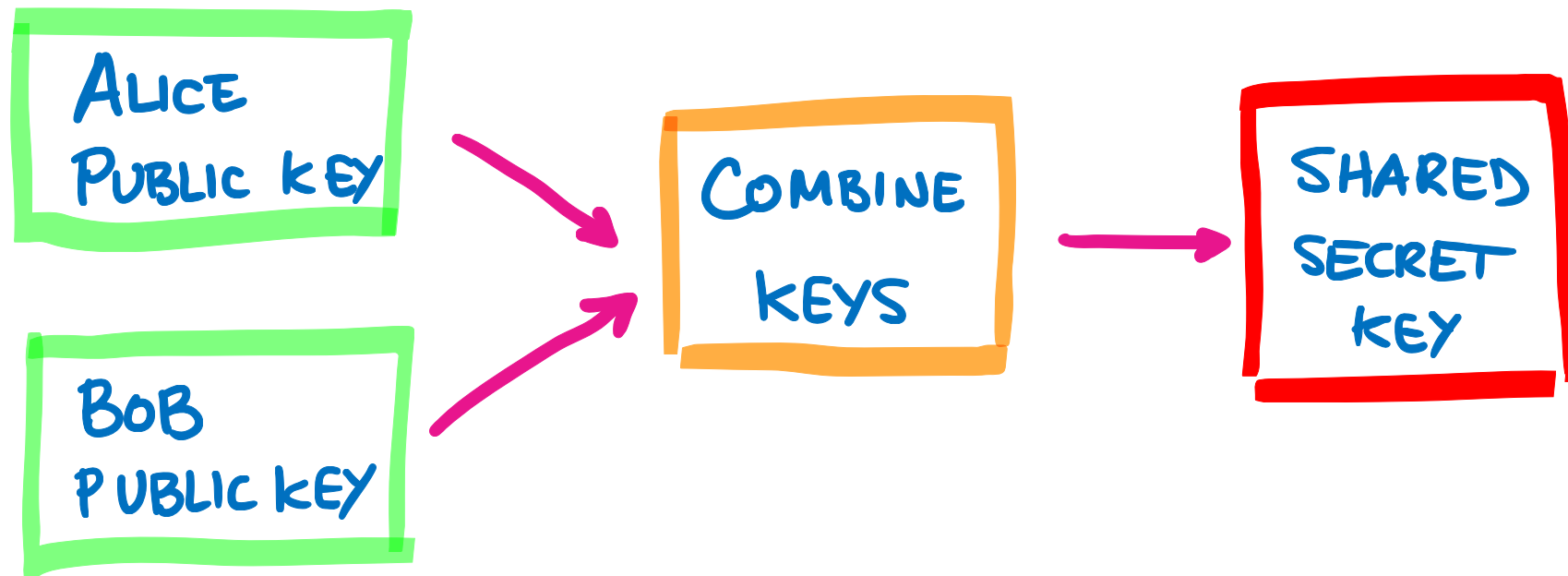
Possible Strings
"Vibrations"



String
Closed Path

Complexity "Sweet Spot"

Public Key Cryptography: Public Key Exchange



DIFFIE - HELLMAN (on $\mathbb{Z}/p\mathbb{Z}$)

1) Agree on a (public) prime p and base g

2) Alice :
• chooses secret exponent a
• public key: $A \equiv g^a \pmod{p}$

3) Bob :
• chooses secret exponent b
• public key: $B \equiv g^b \pmod{p}$

4) Alice and Bob compute secret key:

$$S \equiv B^a \equiv A^b \equiv g^{ab} \pmod{p}$$



STRENGTH:

DISCRETE

LOG PROBLEM

Now replace $\mathbb{Z}/p\mathbb{Z}$ by another group that offers:

- efficiency (key storage/transmission economy)
- greater security

Abstract
Group



ALGEBRAIC
GROUP



ABELIAN
VARIETIES
(dim ≥ 1)



ELLIPTIC
CURVES
(dim = 1)

ELLIPTIC CURVE DIFFIE-HELLMAN

1) Agree on p , ell. curve E/\mathbb{F}_p , and $G \in E(\mathbb{F}_p)$

2) Alice :
· choose a
· publish $A = [a](G) \in E(\mathbb{F}_p)$

3) Bob :
· choose b
· publish $B = [b](G) \in E(\mathbb{F}_p)$

4) Alice and Bob compute secret key:

$$S = [a](B) = [b](A) = [ab](G)$$

EXAMPLE
WhatsApp

Use RSA directly to encrypt?

STRENGTH:

FACTORIZATION OF INTEGERS

ELLIPTIC CURVE FACTORIZATION ALGORITHM!

- Lenstra:
- Pick E/\mathbb{Q} , P of infinite order, compute $P' \equiv P \pmod{n}$
 - Compute $[k](P')$ for k with many small prime divisors p
 - ↳ if it works, new E, P, k
 - ↳ if it doesn't, $\exists plk$ s.t. $p|n$.

ECDSA : Elliptic Curve Digital Signature

↳ EXAMPLES PS3, Bitcoin.



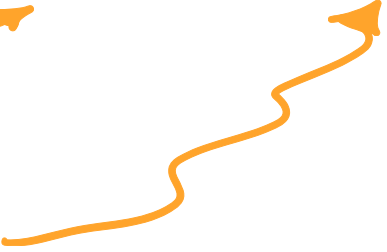

2010: Reused k

2013: alg. for k
caused
"collisions"

- 1) Pick $E, G \in E$ of large prime order p
- 2) Alice:
 - Private key d_A , Public key: $[d_A](G) = A$
 - Message to be signed: z (integer or l -most bits of message)
 - Select "SAFE" $k \in \{1, \dots, n-1\}$
 - Compute $[k](G) = (x_1, y_1) \rightarrow$ SIGNATURE: $(x_1, k^{-1}(z + r \cdot d_A)) \bmod p$
" (r, s) "
- 3) Bob: Verify signature:
 - $A \in E$ ✓
 - $(u_1, u_2) = (z s^{-1}, r s^{-1}) \bmod p$ ✓
 - $(x_1, y_1) = [u_1](G) + [u_2](A)$
if $x_1 \equiv r \bmod p$ ✓

Sweet Spot ?? Maybe just...

LAW OF SMALL NUMBERS

- Dioph. eqn's w/ small degree  Conics, ell. curves
- Riemann surfaces of low genus 
- Abelian varieties of low dimension 
- "Weak" computing power 

Thank You!



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